## Part III: The geometry of lattice congruences on posets of regions

Nathan Reading<br>NC State University

Algebraic and Geometric Combinatorics of Reflection Groups CRM/LaCIM Spring School
UQAM, June 1-2, 2017

Lattice congruences and fans
Cambrian fans
Shards
The shard intersection order

## Main points

$\mathcal{A}$ : simplicial hyperplane arrangement (e.g., Coxeter arrangement).
$\mathcal{F}$ : the fan defined by $\mathcal{A}$.
The poset of regions $\mathcal{P}(\mathcal{A}, B)$ is a polygonal lattice. A congruence $\Theta$ on $\mathcal{P}(\mathcal{A}, B)$ defines a coarsening $\mathcal{F}_{\Theta}$ of $\mathcal{F}$.

When $\mathcal{A}$ is a Coxeter arrangement (so $\mathcal{P}(\mathcal{A}, B)$ is the weak order) and $\Theta$ is a Cambrian congruence $\Theta_{c}, \mathcal{F}_{\Theta}$ is the Cambrian fan $\mathcal{F}_{c}$. It is the normal fan of a realization of the $W$-associahedron (Hohweg-Lange-Thomas) and contains information about $\mathbf{g}$-vectors and $\mathbf{c}$-vectors (R.-Speyer).

A simple geometric condition cuts hyperplanes of $\mathcal{A}$ into shards, which form a geometric model for join-irreducible elements, forcing, and canonical join complex.

## Section III.a: Lattice congruences and fans

## The poset of regions (Edelman, 1985)

$\mathcal{A}$ : a (central) hyperplane arrangement in a real vector space.
Regions: connected components of the complement of $\mathcal{A}$.
$B$ : a distinguished "base" region.
Separating set of a region $R$ :

$$
S(R)=\{\text { hyperplanes of } \mathcal{A} \text { separating } R \text { from } B\}
$$

Poset of regions $\mathcal{P}(\mathcal{A}, B)$ is the set of regions with

$$
Q \leq R \text { if and only if } S(Q) \leq S(R)
$$

Alternately, take the zonotope dual to $\mathcal{A}$ and direct its 1 -skeleton by a linear functional.

Proposition. If $\mathcal{A}$ is a Coxeter arrangement for $W$, then $w \mapsto w B$ is an isomorphism from the weak order on $W$ to $\mathcal{P}(\mathcal{A}, B)$.

## Fans

$V$ : a real vector space.
Closed cone $C \subseteq V$ : closed under nonnegative scaling, addition.
Fan $\mathcal{F}$ : A collection of closed cones such that:
If $C \in \mathcal{F}$ then all faces of $C$ are in $\mathcal{F}$.
If $C, D \in \mathcal{F}$ then $C \cap D$ is a face of $C$ and of $D$.
$\mathcal{F}$ is complete if $\cup \mathcal{F}=V$.
Example. The normal fan of a polytope $P$ in $V$.
Define an equivence relation on functionals in the dual space to $V$ with $f \equiv f^{\prime}$ if and only if $f$ and $f^{\prime}$ are maximized on the same face of $P$. Cones in $\mathcal{F}$ are closures of equivalent classes.

For example, a polygon and its normal fan:


## Coarsening fans by lattice congruences

Every central hyperplane arrangement defines a fan.
(Cones are the regions, together with all their faces.)
This is the normal fan of the corresponding zonotope.
Simplicial fan: all cones are simplicial.
Simplicial hyperplane arrangement: cuts space into a simplicial fan.
Theorem (Bjorner, Edelman, Ziegler, 1987). If $\mathcal{A}$ is simplicial then $\mathcal{P}(\mathcal{A}, B)$ is a lattice for any base region $B$.

Theorem. If $\mathcal{A}$ is simplicial then $\mathcal{P}(\mathcal{A}, B)$ is a polygonal lattice.
For any lattice congruence $\Theta$ on $\mathcal{P}(\mathcal{A}, B)$, define a collection $\mathcal{F}_{\Theta}$ of cones, closed under passing to faces.

Maximal cones of $\mathcal{F}_{\Theta}$ are unions, over congruence classes of $\Theta$, of maximal cones of the fan defined by $\mathcal{A}$.

Theorem (R., 2004). $\mathcal{F}_{\Theta}$ is a fan.

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Example: $\mathcal{F}_{\Theta}$ for a congruence on the weak order on $S_{4}$

$$
\mathcal{F}_{\Theta}=\text { normal fan of associahedron. } \mathcal{P}(\mathcal{A}, B) / \Theta=\text { Tamari lattice. }
$$



Part III: The geometry of lattice congruences on posets of regions


Lattice congruences and fans

## Why $\mathcal{F}_{\Theta}$ is a fan

A complete fan $\mathcal{F}^{\prime}$ coarsens a complete fan $\mathcal{F}$ if each face of $\mathcal{F}^{\prime}$ is a union of faces of $\mathcal{F}$.

Adjacency graph $\mathcal{G}$ of $\mathcal{F}$ :
Vertices are the full-dimensional cones of $\mathcal{F}$
Edges are the pairs of adjacent full-dimensional cones
A fan $\mathcal{F}^{\prime}$ coarsening $\mathcal{F}$ is determined by its edge set: the set of edges connecting adjacent full-dimensional cones of $\mathcal{F}$ that are contained in the same face of $\mathcal{F}^{\prime}$.

One can show that $\mathcal{F}_{\Theta}$ is a fan as a special case of a characterization of which edge sets correspond to fan coarsenings.

The characterization is very general (coarsenings of polytopal complexes), but we'll phrase it for fans coming from hyperplane arrangements.

## Why $\mathcal{F}_{\Theta}$ is a fan (continued)

When $\mathcal{F}$ comes from a hyperplane arrangement, $\mathcal{G}$ is the vertex-edge graph of the zonotope $\mathcal{Z}$ dual to $F$.

The polygon property of a set $\mathcal{E}$ of edges of $\mathcal{G}$ :
For every $2 k$-gonal face $P$ of $\mathcal{Z}$, whenever $\mathcal{E}$ contains any $k-1$ consecutive edges of $P$, then $\mathcal{E}$ also contains the opposite $k-1$ consecutive edges of $P$.

Theorem (R., 2010). Let $\mathcal{Z}$ be a zonotope and let $\mathcal{F}$ be the normal fan of $\mathcal{Z}$. Then a set $\mathcal{E}$ of edges of $\mathcal{Z}$ is the edge set of a fan coarsening $\mathcal{F}$ if and only if $\mathcal{E}$ has the polygon property.

The special case where $\mathcal{Z}$ is the permutohedron is due to Morton, Pachter, Shiu, Sturmfels, and Wienand.

## Why $\mathcal{F}_{\Theta}$ is a fan (continued)

Polygon property: For every $2 k$-gonal face $P$ of $\mathcal{Z}$, whenever $\mathcal{E}$ contains any $k-1$ consecutive edges of $P$, then $\mathcal{E}$ also contains the opposite $k-1$ consecutive edges of $P$.


Theorem (repeated). A set $\mathcal{E}$ of edges of $\mathcal{Z}$ is the edge set of a fan coarsening $\mathcal{F}$ if and only if $\mathcal{E}$ has the polygon property.

Forcing says:


Conclude: If $\mathcal{E}$ is chosen by edge forcing, $\mathcal{F}^{\prime}$ is a fan coarsening $\mathcal{F}$. That is, $\mathcal{F}_{\Theta}$ is a fan coarsening $\mathcal{F}$.

## Recap of Section III.a: Lattice congruences and fans

When $\mathcal{A}$ is simplicial, $\mathcal{P}(\mathcal{A}, B)$ is a polygonal lattice.
$\mathcal{F}$ is the simplicial fan defined by $\mathcal{A}$ (maximal cones are regions).
Given a congruence $\Theta$, for each $\Theta$-class, take the union of the corresponding regions.

These unions are the maximal cones of a complete fan $\mathcal{F}_{\Theta}$ that coarsens $\mathcal{F}$.

## Questions?

## Section III.b: Cambrian fans

## Cambrian fan

The Cambrian fan $\mathcal{F}_{c}$ is $\mathcal{F}_{\Theta_{c}}$ where $\Theta_{c}$ is the Cambrian congruence. That is, maximal cones are unions (over $\Theta_{c}$-classes) of maximal cones of the Coxeter fan.

Theorem (R., Speyer, 2006).
The bijection $\mathrm{cl}_{c}:\{c$-sortables $\} \rightarrow$ \{clusters $\}$ induces a combinatorial isomorphism between the Cambrian fan $\mathcal{F}_{c}$ and the normal fan of the generalized associahedron.

Theorem (Hohlweg, Lange, Thomas, 2010). The Cambrian fan $\mathcal{F}_{c}$ is the normal fan of a realization of the generalized associahedron. (They gave an explicit construction.)

## Example $\left(W=B_{2}, c=s_{1} s_{2}\right)$



The
bijection




## Cambrian fans from the combinatorics of sortable elements

$v$ : a $c$-sortable element of $W$ with $c$-sorting word $a_{1} \cdots a_{k}$.
Recall: $c^{\infty}=s_{1} \cdots s_{n}\left|s_{1} \cdots s_{n}\right| s_{1} \cdots s_{n} \mid \cdots$
For each $s_{i} \in S$, there is a leftmost instance of $s_{i}$ in $c^{\infty}$ which is not in the subword of $c^{\infty}$ corresponding to $a_{1} \cdots a_{k}$.

Let $a_{1} \cdots a_{j}$ be the initial segment of $a_{1} \cdots a_{k}$ consisting of those letters that occur in $c^{\infty}$ before the omission of $s_{i}$. Say $a_{1} \cdots a_{k}$ skips $s_{i}$ after $a_{1} \cdots a_{j}$.
Define: $C_{c}^{s_{i}}(v)=a_{1} \cdots a_{j} \cdot \alpha_{i}$.
$C_{c}(v)=\left\{C_{c}^{s_{i}}: s_{i} \in S\right\}$.
Cone $_{c}(v)=\left\{\mathbf{x} \in\left(\mathbb{R}^{n}\right)^{*}:\left\langle\mathbf{x}, C_{c}^{s_{i}}(v)\right\rangle \geq 0 \forall s_{i} \in S\right\}$.
Theorem (R., Speyer, 2014). The maximal cones of the Cambrian fan $\mathcal{F}_{c}$ are Cone $_{c}(v)$ as $v$ runs over all $c$-sortable elements.

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} \cdots \\
c \text {-sortable: } & 1, \quad s_{1}, & s_{1} s_{2}, \quad s_{1} s_{2} \mid s_{1}, \\
\text { not } c \text {-sortable: } & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip | $C_{c}^{s_{i}(v)}$ |
| :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ |  |  |
| $s_{1}$ | $s_{2}$ |  |  |
| $s_{1} s_{2}$ | $s_{1}$ |  |  |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |
| $s_{2}$ | $s_{2}$ |  |  |
|  | $s_{1}$ |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} \mid s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{2} \mid s_{1} s_{2} & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced $\quad\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
| $s_{1}$ | $s_{2}$ |  |  |
| $s_{1} s_{2}+\alpha_{2}$ |  |  |  |
|  | $s_{2}$ |  |  |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |
| $s_{2}$ | $s_{2}$ |  |  |
|  | $s_{1}$ |  |  |
|  | $s_{2}$ |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
s_{1} s_{2}, & s_{1} s_{2} \mid s_{1}, & s_{1} s_{2} \mid s_{1} s_{2}, \\
\text { not } c \text {-sortable: } & s_{2} \mid s_{1}, & s_{2} \mid s_{1} s_{2}
\end{array}
$$

$\left.\begin{array}{l|l|l|l}v & s_{i} & \text { skip } & C_{c}^{s_{i}}(v) \\ \hline 1 & s_{1} & \text { unforced } & \left(s_{1} \text { reduced }\right) \\ & s_{2} & \text { unforced } & \left(s_{2} \text { reduced }\right)\end{array}\right)$

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} \mid s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{2} \mid s_{1} s_{2} & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ |  |  |  |
| $s_{1} s_{2}$ | $s_{1}$ |  |  |  |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |  |
| $s_{2}$ | $s_{2}$ |  |  |  |
|  | $s_{2}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} \mid s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{1}\left|s_{1} s_{2}\right| s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |  |
| $s_{2}$ | $s_{2}$ |  |  |  |
|  | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |
|  |  |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{1}, & s_{2} s_{2} \mid s_{1} s_{1} & s_{1} s_{2},
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ |  |  |  |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |  |
| $s_{2}$ | $s_{2}$ |  |  |  |
|  | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{1}, & s_{2} \mid s_{1} s_{2} s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |  |
| $s_{2}$ | $s_{2}$ |  |  |  |
|  | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} \cdots$
$c$-sortable: $\quad 1, s_{1}, s_{1} s_{2}, s_{1} s_{2}\left|s_{1}, s_{1} s_{2}\right| s_{1} s_{2}, s_{2}$ not $c$-sortable: $\quad s_{2}\left|s_{1}, s_{2}\right| s_{1} s_{2}$


## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} \cdots$
$c$-sortable: $\quad 1, s_{1}, s_{1} s_{2}, s_{1} s_{2}\left|s_{1}, s_{1} s_{2}\right| s_{1} s_{2}, s_{2}$ not $c$-sortable: $\quad s_{2}\left|s_{1}, s_{2}\right| s_{1} s_{2}$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}-\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2} s_{1} s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |
| $s_{2}$ | $s_{1}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} \cdots$
$c$-sortable: $\quad 1, s_{1}, s_{1} s_{2}, s_{1} s_{2}\left|s_{1}, s_{1} s_{2}\right| s_{1} s_{2}, s_{2}$ not $c$-sortable: $\quad s_{2}\left|s_{1}, s_{2}\right| s_{1} s_{2}$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}-\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2} s_{1} s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ |  |  |  |
| $s_{2}$ | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{2} \mid s_{1} s_{2} & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}-\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2} s_{1} s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-\alpha_{2}$ |
| $s_{2}$ | $s_{1}$ |  |  |  |
|  | $s_{2}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{2} \mid s_{1} s_{2} & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}-\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2} s_{1} s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-\alpha_{2}$ |
| $s_{2}$ | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ |  |  |  |

## Skips example: $W=B_{2}, \quad c=s_{1} s_{2}$

$$
\begin{array}{lll}
c^{\infty}=s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} s_{1} s_{2} & \cdots \\
c \text {-sortable: } & 1, & s_{1}, \\
\text { not } c \text {-sortable: } & s_{2} s_{2}, & s_{1} s_{2} \mid s_{1}, \\
s_{2} \mid s_{1} s_{2} & s_{1} s_{2} \mid s_{1} s_{2}, & s_{2}
\end{array}
$$

| $v$ | $s_{i}$ | skip |  | $C_{c}^{s_{i}}(v)$ |
| :--- | :--- | :--- | :--- | :--- |
| 1 | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2}\right.$ reduced $)$ | $2 \alpha_{1}+\alpha_{2}$ |
| $s_{1} s_{2}$ | $s_{1}$ | unforced | $\left(s_{1} s_{2} s_{1}\right.$ reduced $)$ | $\alpha_{1}+\alpha_{2}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-2 \alpha_{1}-\alpha_{2}$ |
| $s_{1} s_{2} s_{1}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}-\alpha_{2}$ |
|  | $s_{2}$ | unforced | $\left(s_{1} s_{2} s_{1} s_{2}\right.$ reduced $)$ | $\alpha_{2}$ |
| $s_{1} s_{2} s_{1} s_{2}$ | $s_{1}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{1}\right.$ not reduced $)$ | $-\alpha_{1}$ |
|  | $s_{2}$ | forced | $\left(s_{1} s_{2} s_{1} s_{2} s_{2}\right.$ not reduced $)$ | $-\alpha_{2}$ |
| $s_{2}$ | $s_{1}$ | unforced | $\left(s_{1}\right.$ reduced $)$ | $\alpha_{1}$ |
|  | $s_{2}$ | forced | $\left(s_{2} s_{2}\right.$ not reduced $)$ | $-\alpha_{2}$ |

## The Cambrian fan by skips

$$
W=B_{2}, c=s_{1} s_{2}
$$

$C_{c}(v)$ shown in red.


## The Cambrian fan by skips

$$
W=B_{2}, c=s_{1} s_{2}
$$

$C_{c}(v)$ shown in red.


## Recap of Section III.b: Cambrian fans

The Cambrian fan is $\mathcal{F}_{c}=\mathcal{F}_{\Theta_{c}}$ for $\Theta_{c}$ the $c$-Cambrian congruence.
It is the normal fan of a generalized associahedron.
Its geometry can be read off from the combinatorics of $c$-sortable elements (skips in c-sorting words).

Questions?

Section III.c: Shards

## Recall: Combinatorial models

When we talked about noncrossing arc diagrams, we said we wanted a combinatorial model for congruences. Specifically, we wanted a set of objects

- in bijection with join-irreducible elements of $W$.
- with a compatibility relation modeling edges of the CJC (so pairwise compatible sets of objects are in bijection with $W$ ).
- with forcing among j.i. elements read off combinatorially.


## Recall: Combinatorial

When we talked about noncrossing arc diagrams, we said we wanted a combinatorial model for congruences. Specifically, we wanted a set of objects

- in bijection with join-irreducible elements of $W$.
- with a compatibility relation modeling edges of the CJC (so pairwise compatible sets of objects are in bijection with $W$ ).
- with forcing among j.i. elements read off combinatorially.

Now we'll consider a general geometric model based on shards.
(It's not so un-combinatorial... In some sense noncrossing diagrams are shards in type A.)

## What shards are

To make the fan $\mathcal{F}_{\Theta}$ for a congruence $\Theta$ on the weak order, we glue cones of the Coxeter fan together according to congruence classes.

So: contracting an edge means removing the wall between two adjacent cones.

A shard is (the union of) a maximal collections of walls which must always be removed together in a lattice congruence. Each shard turns out to consist of walls all in the same hyperplane.

## Example.




We describe a congruence by specifying which shards are removed. Edge-forcing also implies some forcing relations among shards.

## What shards are

To make the fan $\mathcal{F}_{\Theta}$ for a congruence $\Theta$ on the weak order, we glue cones of the Coxeter fan together according to congruence classes.

So: contracting an edge means removing the wall between two adjacent cones.

A shard is (the union of) a maximal collections of walls which must always be removed together in a lattice congruence. Each shard turns out to consist of walls all in the same hyperplane.*

Example.


*Hypotheses: Weak order on a finite Coxeter group or congruence uniform poset of regions of a simplicial arrangement.

We describe a congruence by specifying which shards are removed. Edge-forcing also implies some forcing relations among shards.


Part III: The geometry of lattice congruences on posets of regions
Shards



Part III: The geometry of lattice congruences on posets of regions
Shards












Part III: The geometry of lattice congruences on posets of regions
Shards


Part III: The geometry of lattice congruences on posets of regions
Shards







Part III: The geometry of lattice congruences on posets of regions
Shards




Part III: The geometry of lattice congruences on posets of regions
Shards


Part III: The geometry of lattice congruences on posets of regions
Shards


## Shards, defined purely geometrically

Shards in a dihedral (or "rank 2") Coxeter group: The two hyperplanes bounding the "identity region" are not cut. The remaining hyperplanes are cut in half.


Important technical point: all of the shards contain the origin. We "cut" along the intersection of the hyperplanes, then take closures of the pieces.

## Shards, defined purely geometrically

Shards in a dihedral (or "rank 2") Coxeter group: The two hyperplanes bounding the "identity region" are not cut. The remaining hyperplanes are cut in half.


Important technical point: all of the shards contain the origin. We "cut" along the intersection of the hyperplanes, then take closures of the pieces. Technical point about to disappear...

## Shards, defined purely geometrically

Shards in a dihedral (or "rank 2") Coxeter group: The two hyperplanes bounding the "identity region" are not cut. The remaining hyperplanes are cut in half.


In higher ranks, we do this cutting in every rank-2 subarrangement.
Why is this the same as the definition by "maximal collections of walls which must always be removed together?" Because the lattice is polygonal!

## Shards, defined purely geometrically

Shards in a dihedral (or "rank 2") Coxeter group: The two hyperplanes bounding the "identity region" are not cut. The remaining hyperplanes are cut in half.


In higher ranks, we do this cutting in every rank-2 subarrangement.
Why is this the same as the definition by "maximal collections of walls which must always be removed together?" Because the lattice is polygonal!

Again, need weak order or the congruence uniform simplicial case for these two definitions to coincide.

## Shards in $S_{4}$

## Shards, join-irreducible congruences, and j.i. elements

Shards are in bijection with join-irreducible congruences. (Maximal collections of walls which must always be removed together are maximal collections of edges that must be contracted together!)

Again, this is for weak order or congruence uniform simplicial posets of regions. So also shards are in bijection with j.i. elements.
In the non congruence uniform case, the bijection from shards to j.i. elements still works, taking the geometric definition of shards.
In any case, the bijection sends a shard to the lowest region above that shard.


Shards

## Compatibility of shards

Two shards are compatible (i.e. form an edge in the canonical join complex) if and only if their relative interiors intersect.


We know that the CJC is flag. Therefore, faces of the CJC are sets of shards that pairwise intersect in their relative interiors.

## Forcing among shards

(i.e. removing one shard forces removal of others)

Facets of shards are maximal proper faces of the shards (codimension 2 in the ambient space).

The shard digraph: $\Sigma_{1} \rightarrow \Sigma_{2}$ iff $\Sigma_{1}$ has a codimension-2 (in the ambient space) intersection with a facet of $\Sigma_{2}$. This can only happen if $\Sigma_{1}$ is in a hyperplane that "cuts" $\Sigma_{2}$ to create that facet.
A shard $\Sigma$ forces another shard $\Sigma^{\prime}$ if and only if there is a directed path from $\Sigma$ to $\Sigma^{\prime}$ in the shard digraph. (Again, because the lattice is polygonal!)

## Forcing among shards

(i.e. removing one shard forces removal of others)

Facets of shards are maximal proper faces of the shards (codimension 2 in the ambient space).

The shard digraph: $\Sigma_{1} \rightarrow \Sigma_{2}$ iff $\Sigma_{1}$ has a codimension-2 (in the ambient space) intersection with a facet of $\Sigma_{2}$. This can only happen if $\Sigma_{1}$ is in a hyperplane that "cuts" $\Sigma_{2}$ to create that facet.

A shard $\Sigma$ forces another shard $\Sigma^{\prime}$ if and only if there is a directed path from $\Sigma$ to $\Sigma^{\prime}$ in the shard digraph. (Again, because the lattice is polygonal!)

Some simplicial hyperplane arrangements have non congruence uniform posets of regions: These are the cases where the shard digraph has oriented cycles.

The proof that the weak order is congruence uniform consists of showing that its shard digraph has no directed cycles.

## Shard removal, forcing and fans in $S_{4}$



## Shard removal, forcing and fans in $S_{4}$



## Shard removal, forcing and fans in $S_{4}$



## Shard removal, forcing and fans in $S_{4}$



## Shard removal, forcing and fans in $S_{4}$



## Recap of Section III.c: Shards

Shards are pieces of hyperplanes that constitute a model for join-irreducible elements and forcing in simplicial hyperplane arrangements.

Compatibility of shards means intersecting in their relative interiors.

Forcing is described in terms of incidence relations among shards.

Questions?

## Section III.d: The shard intersection order

## The shard intersection order

Initial motivation: The lattice property for the noncrossing partition lattice was first proved uniformly by Brady and Watt (2005), and differently (for $W$ crystallographic) Ingalls and Thomas (2006).

Shard intersections give a new proof that $\mathrm{NC}(W)$ is a lattice: Construct a lattice ( $W, \preceq$ ) on the elements of $W$, and identify a sublattice of $(W, \preceq)$ isomorphic to $\mathrm{NC}(W)$.

Beyond the initial motivation: ( $W, \preceq$ ) turns out to have very interesting properties, analogous to the properties of $\mathrm{NC}(W)$.
Proofs are simple and natural in the Coxeter context. (More broadly: in the context of simplicial hyperplane arrangements.)

This approach brings to light how $\mathrm{NC}(W)$ arises naturally in the context of semi-invariants of quivers.

There are intriguing connections to certain "pulling" triangulations of associahedra and permutohedra.

## The shard intersection order

Let $\Psi(W)$ be the set of arbitrary intersections of shards.
We partially order this set by reverse containment.
Immediate: $(\Psi(W), \supseteq)$ is a join semilattice. (Join is intersection.)
It also has a unique minimal element (the empty intersection, i.e. the ambient vector space), so it is a lattice. Also immediate: $(\Psi(W), \supseteq)$ is atomic.
Less obvious: $(\Psi(W), \supseteq)$ is graded (ranked by codimension) and coatomic.

Surprising: The elements of $\Psi(W)$ are in bijection with the elements of $W$.

$$
w \in W \longleftrightarrow \text { a region } R \longleftrightarrow \bigcap\{\text { shards below } R\}
$$

In particular, $(\Psi(W), \supseteq)$ induces a partial order $\preceq$ on $W$.
Also surprising: Every lower interval in $(\Psi(W), \supseteq)$ is isomorphic to $\left(\Psi\left(W_{J}\right), \supseteq\right)$ for some standard parabolic subgroup $W_{J}$.

## Shard intersections in $I_{2}(5)$

The poset $\left(\Psi\left(I_{2}(5)\right), \supseteq\right)$ has $\mathbb{R}^{2}$ as its unique minimal element and the origin as its unique maximal element. The 8 (1-dimensional) shards are pairwise incomparable under containment, and live at rank 1 (i.e. codimension 1).

The poset $\left(I_{2}(5), \preceq\right)$ has 1 as its unique minimal element and $w_{0}$ as its unique maximal element. The other 8 elements of $W$ are pairwise incomparable and live at at rank 1 .

## Shard intersections in $S_{4}$



## The shard intersection lattice on $S_{4}$



## Shard intersections and lattice congruences

Since shards are so central to lattice congruences on the weak order, it is perhaps not surprising that lattice congruences "play nicely" with the shard intersection order.

Specifically, let $\pi_{\downarrow}^{\Theta}(W)$ be the collection of "bottom elements" of congruences classes of a congruence $\Theta$. Then the restriction $\left(\pi_{\downarrow}^{\Theta}(W), \preceq\right)$ is a lattice and a join-sublattice of $\left.(W, \preceq)\right)$.

For $\Theta_{c}$ the $c$-Cambrian congruence, the lattice $\left(\pi_{\downarrow}{ }_{\downarrow}, \preceq\right)$-the restriction to $c$-sortable elements-is isomorphic to $\mathrm{NC}_{c}(W)$.

As a consequence, $\mathrm{NC}_{c}(W)$ is a lattice. (In fact, $\mathrm{NC}_{c}(W)$ is a sublattice of $(W, \preceq)$.)

The earlier proof (by Brady and Watt) that $\mathrm{NC}(W)$ is a lattice also used the polyhedral geometry of cones. Their proof is "dual" to the new proof (in the broadest outlines but not in any of the details).






## Properties of ( $W, \preceq$ ) and NC( $W$ )

| $(W, \preceq)$ | $\mathrm{NC}(W)$ |
| :--- | :--- |
| Lattice | Lattice (sublattice of $(W, \preceq))$ |
| Weaker than weak order | Weaker than Cambrian lattice |
| Atomic and coatomic | Atomic and coatomic |
| Graded ( $W$-Eulerian numbers) | Graded ( $W$-Narayana) |
| Not self-dual | Self-dual |
| Lower intervals $\cong\left(W_{J}, \preceq\right)$ | Lower intervals $\cong \mathrm{NC}\left(W_{J}\right)$ |
| Möbius number: $\pm$ number of <br> "positive" elements of $W$. | Möbius number: $\pm$ number of <br> "positive" elements of $\mathrm{NC}(W)$. |

## Details on the Möbius number

Theorem. The Möbius function of ( $W, \preceq$ ) satisfies

$$
\mu\left(1, w_{0}\right)=\sum_{J \subseteq S}(-1)^{|J|}\left|W_{J}\right| .
$$

Proof. Since lower intervals $[1, w]$ are isomorphic to $\left(W_{\operatorname{Des}(w)}, \preceq\right)$, checking the defining recursion for $\mu$ becomes

$$
\sum_{w \in W} \sum_{J \subseteq \operatorname{Des}(w)}(-1)^{|J|}\left|W_{J}\right|=\sum_{J \subseteq S}(-1)^{|J|}\left|W_{J}\right| \sum_{\substack{w \in W_{\text {s.t. }} \subseteq \subseteq \operatorname{Des}(w)}} 1 .
$$

The inner sum is $|W| /\left|W_{J}\right|$, the number of maximal-length representatives of cosets of $W_{J}$ in $W$. Thus the double sum reduces to zero.

## Properties of ( $W, \preceq$ ) and $\mathrm{NC}(W)$ (continued)

| $(W, \preceq)$ | $\mathrm{NC}(W)$ |
| :--- | :--- |
| Recursion counting maximal <br> chains: sum over max'l proper <br> standard parabolic subgroups. | Recursion counting maximal <br> chains: sum over max'l proper <br> standard parabolic subgroups. |
| $\mathrm{MC}(W)=$ | (R., 2007.) |
| $\sum_{s \in S}\left(\frac{\|W\|}{\left\|W_{\langle s\rangle}\right\|}-1\right) \mathrm{MC}\left(W_{\langle s\rangle}\right)$ | $\mathrm{MC}(W)=\frac{h}{2} \sum_{s \in S} \mathrm{MC}\left(W_{\langle s\rangle}\right)$. |

These types of recursions are very natural in the context of Coxeter groups/root systems. For example:

1. Recursions for the $W$-Catalan number.
2. Volume of $W$-permutohedron (weight polytope). This follows from Postnikov's formula in terms of $\Phi$-trees.

## Properties of $(W, \preceq)$ and $N C(W)$ (concluded)

| $(W, \preceq)$ | $\mathrm{NC}(W)$ |
| :--- | :--- |
| Maximal chains $\longleftrightarrow$ maximal <br> simplices in a pulling triangula- <br> tion of the $W$-permutohedron. | Maximal chains $\longleftrightarrow$ maximal <br> simplices in a pulling triangu- <br> lation of the $W$-associahedron. |
| $\left(S_{n}\right.$ case: Loday described the | $\left(S_{n}\right.$ case: Loday, 2005.) |
| triangulation, 2005.) | (General case: R., 2008.) |
| $k$-Chains $\longleftrightarrow k$-simplices in | $k$-Chains $\longleftrightarrow k$-simplices in |
| the same triangulation of the | the same triangulation of the |
| $W$-permutohedron. | $W$-associahedron. |
|  | $(R ., 2008)$. |

Loday: Noticed that maximal simplices in a certain pulling triangulation of the $S_{n}$-associahedron biject with parking functions. Constructed the analogous triangulation of the $S_{n}$-permutohedron and asked what played the role of parking functions.

## Details on the triangulations

The bijection between intersections of shards and elements of $W$ extends to a bijection between $k$-chains in ( $W, \preceq$ ) and $k$-simplices in a pulling triangulation of the $W$-permutohedron.

In particular: The order complex of ( $W, \preceq$ ) has $f$-vector equal to the $f$-vector of a pulling triangulation of the $W$-permutohedron.

Key point: For any $w \in W$, the lower interval $[1, w]$ in $(W, \preceq)$ is isomorphic to $\left(W_{J}, \preceq\right)$ for some $W_{J}$. The elements of $W_{J}$ are in bijection with vertices of the face below $w$ in the permutohedron.

All of this works for $\mathrm{NC}(W)$ and the $W$-associahedron as well. Maximal chains in $\mathrm{NC}\left(S_{n}\right)$ are in bijection with parking functions, so we recover the Loday result as a special case.

## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example

321


## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## $S_{3}$ Permutohedron example



## Recap of Section III.d: The shard intersection order

Elements of $W$ are in bijection with intersections of shards.

Intersections of shards form a lattice, which can be interpreted as a new lattice structure on $W$.
c-Sortable elements induce a sublattice isomorphic to the lattice of noncrossing partitions.

NC partitions and the shard intersection order have analogous structural and enumerative properties.

Questions?

## References

C. Hohlweg, C. Lange and H. Thomas, Permutahedra and generalized associahedra. Adv. Math., 2011.
N. Reading, Lattice congruences, fans and Hopf algebras. JCTA, 2005.
N. Reading, Noncrossing partitions and the shard intersection order. JACo (2011).
N. Reading and D. E. Speyer, Cambrian fans. JEMS, 2009.
N. Reading and D. E. Speyer, Combinatorial frameworks for cluster algebras. IMRN, 2016.

